Milestone 2: Algorithm Design

# **Milestone 1: Understanding the problems**

**GitHub Repository Link:** <https://github.com/Angellsh/COP4533-Group-13.git>

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**Member Roles:**

* **Jacob Ramos**: Problem 1 & 2 completed both tasks (100% progress)
* **Loubna Benchakouk**: Problem 3 completed (100% progress)
* **Anhelina Liashynska**: GitHub setup
* **Dani Brown**: Gantt Chart
* Roles for upcoming milestones will be assigned later

**Communication methods:** We use Discord for regular communication and Google Docs for document collaboration.

**Project Gantt Chart:**

# **Problem 1**

We are given a matrix of stock prices where each row represents a different stock and each column will represent a different day. We calculate the maximum potential profit for each specific stock using a 1-based index.

Each stock/day combination maximum profit: (1,2,5,15) (2,1,3,9) (3,1,2,2) (4,2,5,7)

Answer: Stock/Day Combination Maximum Profit: (1,2,5,15)

Explanation:

* Stock 1 yields the maximum when bought on the 2nd day and sold on the 5th day for a profit of 15.
* Stock 2 yields the maximum profit when bought on day 1 and sold on day 3 yielding a profit of 9.
* Stock 3 yields the maximum potential profit when bought on day 1 and sold on day 2 yielding a profit of 2.
* Finally, stock 4 yields the maximum potential profit when bought on day 2 and sold on day 5 yielding a profit of 7.

# **Problem 2**

We are given a matrix where each row represents a different stock and each column will represent a different day. We are given an integer k which will represent the maximum number of non-overlapping transactions permitted, in this case k = 3. For each transaction we must buy and sell one stock.

Answer: (4,1,2), (2,2,3), (1,3,5) total profit = 90

Explanation:

1. Stock 4: Buy on the 1st day at price 5, sell on the 2nd day at price 50 for a profit of 45.
2. Stock 2: Buy on the 2nd day at price 20, sell on the 3rd day at price 30 for a profit of 10.
3. Stock 1: Buy on the 3rd day at price 15, sell on the 5th day at price 50 for a profit of 35.
4. Total profit = 45 + 10 + 35 = 90

# **Problem 3**

**Problem Statement**

We are given a matrix where each row represents a different stock and each column will represent a different day. Additionally, we are given an integer c which will represent a cooldown period where we cannot buy any stock for c days after selling any stock. If a stock is sold on day i, the next stock will not be eligible for purchase until day i + c + 1. For this example, c = 2.

Answer: (3,1,3), (3,6,7) total profit = 4 + 7 = 11

Explanation:

1. First transaction we buy stock 3 on day 1 and sell on day 3 for a profit of 4
2. Since the stock was sold on day 3 we cannot purchase another stock till day 6
3. On day 6, we buy stock 3 again and sell on day 7 for a profit of 7
4. The total profit is 11

**Transaction rules:**

1. We can only buy before we sell, and only once per transaction.

2. Resting period: after we sell on day j2 we need to wait until (j2+c+1) day to buy.

3. We can perform multiple transactions on any stock while following the cooldown rule.

4. Main objective is to maximize the total profit across all valid transactions.

**Input:**

We have a matrix A where each:

Row = one stock

Column = one day

A[i][j] = price of stock(i + 1) on day(j + 1)

Matrix A:

|  |  |  |  |  |  |  |  |
| --- | --- | --- | --- | --- | --- | --- | --- |
| Day | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
| Stock\_1 | 7 | 1 | 5 | 3 | 6 | 8 | 9 |
| Stock\_2 | 2 | 4 | 3 | 7 | 9 | 1 | 8 |
| Stock\_3 | 5 | 8 | 9 | 1 | 2 | 3 | 10 |
| Stock\_4 | 9 | 3 | 4 | 8 | 7 | 4 | 1 |
| Stock\_5 | 3 | 1 | 5 | 8 | 9 | 6 | 4 |

Cooldown period: c = 2  
  
To solve this problem we need to find all profitable transactions for each stock(row in the matrix)

1. Choose a buy day and then try all sell days that come after that buy day
2. For each(buy, sell) day, check if the price on the sell day is higher than the price on the buy day.
3. Keep just the profitable pairs(i, j, l)

**Step 1: Identify All Possible Profitable Transactions**

For each stock, we need to check all (buy, sell) pairs where buyDay < sellDay and profit > 0:

**Stock 1: [7, 1, 5, 3, 6, 8, 9]**

|  |  |  |  |  |  |  |
| --- | --- | --- | --- | --- | --- | --- |
| BuyDay | SellDay | BuyPrice | SellPrice | Profit | NextBuy | ValidTransaction |
| 1 | 2 | 7 | 1 | -6 |  |  |
| 3 | 7 | 5 | -2 |  |  |
| 4 | 7 | 3 | -4 |  |  |
| 5 | 7 | 6 | -1 |  |  |
| 6 | 7 | 8 | 1 | Day9(6+2+1) | No |
| 7 | 7 | 9 | 2 | Day10(7+2+1) | No |
| 2 | 3 | 1 | 5 | 4 | Day6(3+2+1) | (6,7) |
| 4 | 1 | 3 | 2 | Day7(4+2+1) | (7,7) |
| 5 | 1 | 6 | 5 | Day8(5+2+1) | No |
| 6 | 1 | 8 | 7 | Day9(6+2+1) | No |
| 7 | 1 | 9 | 8 | Day10(7+2+1) | No |
| 3 | 4 | 5 | 3 | -2 |  |  |
| 5 | 5 | 6 | 1 | Day8(5+2+1) | No |
| 6 | 5 | 8 | 3 | Day9 | No |
| 7 | 5 | 9 | 4 | Day10 | No |
| 4 | 5 | 3 | 6 | 3 | Day8 | No |
| 6 | 3 | 8 | 5 | Day9 | No |
| 7 | 3 | 9 | 6 | Day10 | No |
| 5 | 6 | 6 | 8 | 2 | Day9 | No |
| 7 | 6 | 9 | 3 | Day10 | No |
| 6 | 7 | 8 | 9 | 1 | Day10 | No |

From the table we see that the best combination for Stock 1: (2,7) with profit = 8

**Stock 2: [2, 4, 3, 7, 9, 1, 8]**

|  |  |  |  |  |  |  |
| --- | --- | --- | --- | --- | --- | --- |
| BuyDay | SellDay | BuyPrice | SellPrice | Profit | NextBuy | ValidTransaction |
| 1 | 2 | 2 | 4 | 2 | Day5(2+2+1) | (5, 6); (5, 7); (6, 7) |
| 3 |  | 3 | 1 | Day6 | (6, 7) |
| 4 |  | 7 | 5 | Day7 | (7, 7) |
| 5 |  | 9 | 7 | Day8 | No |
| 6 |  | 1 | -1 |  |  |
| 7 |  | 8 | 6 | Day10 | No |
| 2 | 3 | 4 | 3 | -1 |  |  |
| 4 |  | 7 | 3 | Day7 | (7, 7) |
| 5 |  | 9 | 5 | Day8 | No |
| 6 |  | 1 | -3 |  |  |
| 7 |  | 8 | 4 | Day10 | No |
| 3 | 4 | 3 | 7 | 4 | Day7 | (7, 7) |
| 5 |  | 9 | 6 | Day8 | No |
| 6 |  | 1 | -2 |  |  |
| 7 |  | 8 | 5 | Day10 | No |
| 4 | 5 | 7 | 9 | 2 | Day8 | No |
| 6 |  | 1 | -6 |  |  |
| 7 |  | 8 | 1 | Day10 | No |
| 5 | 6 | 9 | 1 | -8 |  |  |
| 7 |  | 8 | -1 |  |  |
| 6 | 7 | 1 | 8 | 7 | Day9 | No |

Best single transaction for Stock 2: (1,5) with profit = 7

**Stock 3: [5, 8, 9, 1, 2, 3, 10]**

|  |  |  |  |  |  |  |
| --- | --- | --- | --- | --- | --- | --- |
| BuyDay | SellDay | BuyPrice | SellPrice | Profit | NextBuy | ValidTransaction |
| 1 | 2 | 5 | 8 | 3 | Day5 | (5, 6); (5, 7); (6, 7) |
| 3 |  | 9 | 4 | Day6 | (6, 7) |
| 4 |  | 1 | -4 |  |  |
| 5 |  | 2 | -3 |  |  |
| 6 |  | 3 | -2 |  |  |
| 7 |  | 10 | 5 | Day10 | No |
| 2 | 3 | 8 | 9 | 1 | Day6 | (6, 7) |
| 4 |  | 1 | -7 |  |  |
| 5 |  | 2 | -6 |  |  |
| 6 |  | 3 | -5 |  |  |
| 7 |  | 10 | 2 | Day10 | No |
| 3 | 4 | 9 | 1 | -8 |  |  |
| 5 |  | 2 | -7 |  |  |
| 6 |  | 3 | -6 |  |  |
| 7 |  | 10 | 1 | Day10 | No |
| 4 | 5 | 1 | 2 | 1 | Day8 | No |
| 6 |  | 3 | 2 | Day9 | No |
| 7 |  | 10 | 9 | Day10 | No |
| 5 | 6 | 2 | 3 | 1 | Day9 | No |
| 7 |  | 10 | 8 | Day10 | No |
| 6 | 7 | 3 | 10 | 7 | Day10 | No |

Best single transaction for Stock 3: (4,7) with profit = 9

**Stock 4: [9, 3, 4, 8, 7, 4, 1]**

|  |  |  |  |  |  |  |
| --- | --- | --- | --- | --- | --- | --- |
| BuyDay | SellDay | BuyPrice | SellPrice | Profit | NextBuy | ValidTransaction |
| 1 | 2 | 9 | 3 | -6 |  |  |
| 3 |  | 4 | -5 |  |  |
| 4 |  | 8 | -1 |  |  |
| 5 |  | 7 | -2 |  |  |
| 6 |  | 4 | -5 |  |  |
| 7 |  | 1 | -8 |  |  |
| 2 | 3 | 3 | 4 | 1 | Day6(3+2+1) | (6, 7) |
| 4 |  | 8 | 5 | Day7 | (7, 7) |
| 5 |  | 7 | 4 | Day8 | No |
| 6 |  | 4 | 1 | Day9 | No |
| 7 |  | 1 | -2 |  |  |
| 3 | 4 | 4 | 8 | 4 | Day7 | (7, 7) |
| 5 |  | 7 | 3 | Day8 | No |
| 6 |  | 4 | 0 | Day9 | No |
| 7 |  | 1 | -3 |  |  |
| 4 | 5 | 8 | 7 | -1 |  |  |
| 6 |  | 4 | -4 |  |  |
| 7 |  | 1 | -7 |  |  |
| 5 | 6 | 7 | 4 | -3 |  |  |
| 7 |  | 1 | -6 |  |  |
| 6 | 7 | 4 | 1 | -3 |  |  |

Best single transaction for Stock 4: (2,4) with profit = 5

**Stock 5: [3, 1, 5, 8, 9, 6, 4]**

|  |  |  |  |  |  |  |
| --- | --- | --- | --- | --- | --- | --- |
| BuyDay | SellDay | BuyPrice | SellPrice | Profit | NextBuy | ValidTransaction |
| 1 | 2 | 3 | 1 | -2 |  |  |
| 3 |  | 5 | 2 | Day6(3+2+1) | (6, 7) |
| 4 |  | 8 | 5 | Day7 | (7, 7) |
| 5 |  | 9 | 6 | Day8 | No |
| 6 |  | 6 | 3 | Day9 | No |
| 7 |  | 4 | 1 | Day10 | No |
| 2 | 3 | 1 | 5 | 4 | Day6 | (6, 7) |
| 4 |  | 8 | 7 | Day7 | (7, 7) |
| 5 |  | 9 | 8 | Day8 | No |
| 6 |  | 6 | 5 | Day9 | No |
| 7 |  | 4 | 3 | Day10 | No |
| 3 | 4 | 5 | 8 | 3 | Day7 | (7, 7) |
| 5 |  | 9 | 4 | Day8 | No |
| 6 |  | 6 | 1 | Day9 | No |
| 7 |  | 4 | -1 |  |  |
| 4 | 5 | 8 | 9 | 1 | Day8 | No |
| 6 |  | 6 | -2 |  |  |
| 7 |  | 4 | -4 |  |  |
| 5 | 6 | 9 | 6 | -3 |  |  |
| 7 |  | 4 | -5 |  |  |
| 6 | 7 | 6 | 4 | -2 |  |  |

Best single transaction for Stock 5: (2,5) with profit = 8

Since we know the best individual transactions per stock. Now we check if we can combine some of them to build a valid sequence.

Starting with stock 1, the best transaction is: buy on day 2, sell on day 7 with profit = 8. After applying the cooldown rule the next valid buy day is day 10 but our max day is 7. Therefore, we can’t combine it with any other transaction

* Sequence (1, 2, 7) with total profit = 8

Stock 2: The best transaction is to buy on day 1 and sell on day 5 with profit = 7 and since the next buy day is day 8 we can’t make an extra transaction.

* Sequence (2, 1, 5) with total profit = 7

but we have another transaction with a smaller profit of 2 if we buy on day 1 and sell on day 2, after the resting period we can buy stock 3 on day 5, sell on day 7 with profit =8

* Sequence (2, 1, 2), (3, 5, 7) with total profit = 10

Stock 3 we found that the best transaction is to buy on day 4, sell on day 7 with profit = 9 and since we need to wait for day10 (invalid) to make another transaction

* Sequence (3, 4, 7) with total profit = 9

But if we buy on day 1 and sell on day 3 with profit = 4, we can combine it with Stock 2 on day 6 after the cooldown period, we buy on day 6 and sell on day 7 with profit = 7

* Sequence (3, 1, 3), (2, 6, 7) with total profit = 11

Stock 4, we have the best profit = 5 if we buy on day 2 and sell on day 4, since the next valid buy day is day 7 and there is no available transaction starting day 7

* Sequence (4, 2, 4) with total profit = 5

For Stock 5 the best transaction is when we buy on day 2 and sell on day 5 with profit = 8, after applying the cooldown rule, we don’t get a valid day

* Sequence (5, 2, 5) with total profit = 8

From the above, the maximum profit = 11 from the sequence (3, 1, 4), (2, 6, 7)

* To achieve the maximum profit, buy 3rd stock on day 1, sell it on day 3. buy 2nd stock on day 6 and sell it on day 7 adhering to 2 days waiting period

# **Milestone 2: Algorithm Design**

2.1 Task-1: Brute Force Algorithm for Problem1 O(m·n2).

The goal is to find the maximum profit from a single buy/sell transaction on the same stock with one buy before one sell, only one transaction, and return stock index, buy day, sell day, and max profit.

**Assumptions & variable definitions:**

* m: number of stocks (rows in matrix A)
* n: number of days (columns in A)
* A transaction is defined as buying a stock on day j1 and selling it later on day j2, where j1 < j2.
* The transaction must be on the same stock (row).
* The result should be a tuple: (i, j1, j2, profit) where:
* i: index of the chosen stock (1-based index)
* j1: day to buy
* j2: day to sell
* profit = A[i][j2] - A[i][j1]

**Pseudocode:**

Algorithm MaxProfitBruteForce(A, m, n)

Input: matrix A(m × n) representing stock prices

Output: tuple (i, j₁, j₂, profit) representing the best stock and days to buy/sell and max profit

maxProfit ← 0

bestStock ← 0

bestBuyDay ← 0

bestSellDay ← 0

for i ← 0 to m − 1 do

for j₁ ← 0 to n − 2 do

for j₂ ← j₁ + 1 to n − 1 do

profit ← A[i][j₂] − A[i][j₁]

if profit > maxProfit then

maxProfit ← profit

bestStock ← i + 1

bestBuyDay ← j₁ + 1

bestSellDay ← j₂ + 1

if maxProfit = 0 then

return (0, 0, 0, 0) // no profitable transaction

else

return (bestStock, bestBuyDay, bestSellDay, maxProfit)

The algorithm checks all possible pairs of days (j1, j2) for each stock to calculate potential profit, it loops over every stock (m stocks) and for each stock, compares every pair of buy/ sell days.

The algorithm keeps track of the max profit found so far and stores its stock index and days. If no profitable transaction exists (all negative or 0), it returns (0, 0, 0, 0).

2.2 Task-2: Greedy Algorithm forProblem1 O(m·n)

**Assumptions & variable definitions:**

* Only one transaction is allowed per stock
* Buy must occur before sell (j₁ < j₂)
* Each stock is evaluated independently
* Return (0, 0, 0, 0) if no profit is possible

**Pseudocode:**

Algorithm MaxProfitGreedy(A, m, n)

Input: matrix A(m × n) representing stock prices

Output: tuple (i, j₁, j₂, profit) representing the best stock and days to buy/sell and max profit

maxProfit ← 0

bestStock ← 0

bestBuyDay ← 0

bestSellDay ← 0

for i ← 0 to m − 1 do

minPrice ← A[i][0]

minDay ← 0

for j ← 1 to n − 1 do

if A[i][j] − minPrice > maxProfit then

maxProfit ← A[i][j] − minPrice

bestStock ← i + 1

bestBuyDay ← minDay + 1

bestSellDay ← j + 1

if A[i][j] < minPrice then

minPrice ← A[i][j]

minDay ← j

if maxProfit = 0 then

return (0, 0, 0, 0)

else

return (bestStock, bestBuyDay, bestSellDay, maxProfit)

In this greedy version of the algorithm, we optimize the search for the max profit from a single buy/ sell transaction. For each stock, the algorithm keeps the minimum price observed so far (minPrice) and the corresponding day (minDay). As it iterates through the remaining days, it computes the current potential profit by subtracting minPrice from the price on the current day. If this profit exceeds the previously recorded maxProfit, the algorithm updates the optimal transaction details (stock index, buy day, and sell day). If the current day's price is less than minPrice, it becomes the new minPrice. If no profitable transaction exists, the algorithm returns the default tuple (0, 0, 0, 0).

**Time Complexity:** O(m · n)

The Algorithm iterates over each of the (m) stocks

For each stock, it performs a single pass over (n) daily prices

In each pass, comparisons and updates are done in constant time

* Total = O(m · n)

2.3 Task-3: Dynamic Programming Algorithm for Problem1 O(m·n)

Assumptions & variable definitions:

* input stocks is a 2D list of shape m × n
* the transaction must obey buy\_day < sell\_day.

Definitions:

* m – number of stocks
* d – number of days
* opt[m][d] – 2d array with max profit for stocks 1..m up to day d.
* best\_stock – index of stock with the highest profit
* stock\_max – highest profit across all stocks
* min\_ind - The index i such that stock m has its **lowest price so far** at day i, for the current day j, where j > i.

Pseudocode:

**function TASK3(stocks): // stocks is a 2D array: m stocks, n days**

**m = length(stocks)**

**d = length(stocks[0])**

**create 2D array opt[m][d], initialized to 0**

**stock\_max = 0**

**min\_opt = -1**

**best\_stock = -1**

**for stock\_index from 0 to m - 1:**

**min\_ind = 0 // index of minimum price so far**

**for day\_index from 1 to d - 1:**

**opt[stock\_index][day\_index] = MAX(**

**opt[stock\_index][day\_index - 1],**

**stocks[stock\_index][day\_index] - stocks[stock\_index][min\_ind]**

**)**

**if stocks[stock\_index][day\_index] < stocks[stock\_index][min\_ind]:**

**min\_ind = day\_index**

**if opt[stock\_index][d - 1] > stock\_max:**

**stock\_max = opt[stock\_index][d - 1]**

**best\_stock = stock\_index**

**min\_opt = min\_ind**

**for i from min\_opt + 1 to d - 1:**

**if stocks[best\_stock][i] - stocks[best\_stock][min\_opt] == stock\_max:**

**return best\_stock, min\_opt, i, stock\_max // stock ID, buy day, sell day, profit**

**return 0, 0, 0, 0 // if no profit found**

**Explanation**

The algorithm iterates over all stocks, processing each row (stock) sequentially. At each step, it attempts to maximize the transaction amount by either reusing a previously computed transaction or by executing a new transaction, which involves subtracting the minimum value so far from the current value.

The minimum index is reset for each row, and any index i within the current stock row can be a candidate minimum for index j, as long as i < j and stocks[i] = argmin(stocks[i:j-1]).

At the end of each row, the algorithm compares the best local profit for the current stock with the global profit. If the local profit is higher, the global profit, best\_stock, and buy\_day are updated accordingly.

Once all stocks have been processed, the values for best\_stock, buy\_day, and profit are already computed. The algorithm then backtracks to identify the corresponding sell\_day and finally returns the sequence: best\_stock, buy\_day, sell\_day, and profit.

2.4 Task-4: Dynamic Programming Algorithm for Problem2 O(m·n2k)

**Assumptions & variable definitions:**

**Pseudocode:**

2.5 Task-5: Dynamic Programming Algorithm for Problem2 O(m·n·k)

**Assumptions & variable definitions:**

* s – number of stocks
* d – number of days
* opt[m][d][k] – 3d array with max profit for stocks 1..m, up to day d, with at most k transactions.
* cmax[kval] – 1d array with a maximum profit from a single transaction for each transaction count k.
* buy[s][kval] – 2d array storing the best buy day for each stock s and each transaction count k.
* transactions – a list of transaction tuples, each containing (stock, buy\_day, sell\_day)

**Pseudocode:**

**function TASK5 (stocks , kval): // stocks is a 2d array, kval is max allowed transactions**

**profit = -1**

**s = length(stocks)**

**d = length(stocks[0])**

**create 3D array opt[s][d][kval], initialized to 0**

**create 1D array cmax[kval], initialized to -INF**

**create 2D array buy[s][kval], initialized to 0**

**for index n from 1 to d-1:**

**for index m from 0 to s-1:**

**for index k from 0 to kval-1: //Index k starts at 0 but represents the 1st transaction**

**//Updates the current maximum profit for transactions k**

**cmax[k] = max(cmax[k], stocks[m][n] - stocks[m][buy[m][k]])**

**//Recurrance relations for different cases**

**if(k==0 and m==0):**

**opt[m][n][k] = cmax[k]**

**else if(k==0):**

**opt[m][n][k] = max( opt[m-1][n][k], cmax[k])**

**else if(m==0):**

**opt[m][n][k] = max( opt[m][n-1][k],**

**opt[m][n-1][k-1] + cmax[k])**

**else:**

**opt[m][n][k] = max( opt[m-1][n][k],**

**opt[m][n-1][k],**

**opt[m][n-1][k-1] + cmax[k])**

**// Update buy day if current price is lower than previous buy day price**

**if k==0 and stocks[m][n] <= stocks[m][buy[m][k]]:**

**buy[m][k] = n**

**else if( stocks[m][n] <= stocks[m][buy[m][k]]) :**

**buy[m][k] = n**

**// Update buy day if current price is lower than previous buy day price and a new transaction has been made**

**else if(opt[m][n][k] == opt[m][n-1][k-1] + cmax[k] ):**

**buy[m][k] = n**

**cmax[k] = 0**

**// Find maximum profit and corresponding transaction count k**

**k=-1**

**profit = 0**

**for index i from 0 to kval-1:**

**if (opt[s-1][d-1][i] > profit):**

**profit = opt[s-1][d-1][i]**

**k = i**

**if k==-1:**

**return (0, 0, 0, 0)**

**sell = length(d)-1**

**m = length(s) -1**

**// Find last sell day (where the optimal values are reused)**

**transactions = []**

**while sell>0:**

**// Same profit without using current stock, move up**

**if m>0 and opt[m][sell][k] == opt[m-1][sell][k]:**

**m-=1**

**// Same profit as on the previous day, move left**

**else if opt[m][sell][k] == opt[m][sell-1][k]:**

**sell-=1**

**else:**

**diff = opt[m][sell][k] - opt[m][sell-1][k-1]**

**gap = INF**

**cstock = -1**

**cbuy = -1**

**for index i from m down to 0:**

**for index j from sell down to 0:**

**if diff == stocks[i][sell] - stocks[i][j] and sell-j < gap:**

**gap= sell-j**

**cstock = i**

**cbuy = j**

**if cstock==-1 :**

**break**

**transactions.append((cstock, cbuy, sell))**

**sell = cbuy**

**k-=1**

**return reverse(transactions)**

**At each step, the algorithm considers three main steps:**

* **Exclude the current transaction**:

This occurs when including the current transaction does not improve the result. For example, in a strictly decreasing price sequence, it's better to reuse the previous best value from day n-1. In this case, the buy[m][k] (buy day) and cmax (current max) still can reuse the values from the range that is not covered by n-1, unless the current value is the smallest (buy day updated) or the current value is largest (cmax updated).

* **Exclude the current transaction and the current stock:**

This case occurs when skipping the current stock (row m, opt[m][n-1]) and the current transaction at day n produces a better result. In other words, the result from the previous stock(s) at the same day – opt[m-1][n] is better. This situation arises when earlier stocks (rows 0 to m-1) had transactions with a higher potential profit at day n, compared to what the current stock can offer with and without the newly computed transaction across all stocks up to this point (0 to m rows).

* **Include the result from the current stock + the maximum profit** for the available interval

In this case, we compute the result as:  
**opt[m][n-1][k-1] + cmax**, where cmax is the maximum profit that can be obtained across all m stocks up to this point for the available interval [i..j], where i is a sell day for opt[m][n-1][k-1]. Since we intend to include a new transaction, we must use the result from **at most** k-1 **transactions** up to this point (opt[m][n-1][k-1]), and then add the best profit possible for the k-th transaction (cmax).

Because we iterate over the nth day across all stocks (rows) in parallel, it is guaranteed that cmax will not overlap with the previous transaction result stored in opt[n-1] for a given k. This is ensured by two properties of the algorithm:

* The maximum value (cmax) is reset each time a new transaction is added to k, preventing overlap between transactions.
* The dynamic nature of the algorithm ensures that if a higher-profit transaction exists (one that covers a wider valid interval), it will have already been covered in the value of opt[n-1][k-1].

These properties guarantee correctness of the algorithm.

The algorithm iterates over all days n for all stocks m, computing a unique optimal profit value for each transaction count k, resulting in a time complexity of O(mnk). Backtracking is also performed to recover the transaction sequence, wth a complexity of (mn^2) ~ n^3.

Therefore, the complexity of the algorithm is O(mnk) ~ O(n^3)

**Backtracking:**

There are three cases:

* **opt[m][n][k] == opt[m][n-1][k]** if the previous result is reused. In this case, we must decrement the sell variable, as no profit was made on the current day, and the profit was reused without making a transaction.
* **opt[m][n][k] == opt[m-1][n][k]** if the previous stock's result is reused. In this case, we must decrement the stock variable, as the current stock produced at most as good a result as the previous stock.
* If none of the above conditions are true, then we must have reached a point where a **transaction occurred**. By the property of our recurrence relation, we know that we must have computed opt[m][n-1][k-1] with the highest profit for the available interval [i..j], where j is greater than or equal to the sell day for opt[m][n-1][k].

However, we do not store any list to recover this information directly. So, we must iterate over all rows of the stocks array and find the **minimum range** [i..j] that produces the same difference between opt[m][n][k] and opt[m][n-1][k-1].

By finding the smallest range, we can guarantee that there is no better way to compose the current optimal result. And even if there exists a different way to arrive at the same difference between the current optimal and the previous (n-1) optimal, the result obtained is still valid or better.

We exhaustively repeat these steps to recover the entire sequence.

In the end, we **reverse the list** to preserve the original order of transactions.

2.6 Task-6: Dynamic Programming Algorithm for Problem3 O(m·n2)

**Assumptions & variable definitions:**

**Pseudocode:**

2.7 Task-7: Dynamic Programming Algorithm for Problem3 O(m·n)

**Assumptions & variable definitions:**

**Pseudocode:**